

# TWO EXTENSIONS OF LEWIS' S3 WITH PEIRCE LAW†

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ABSTRACT: We define two extensions of Lewis' S3 with two versions of Peirce's Law. We prove that both of them have the Ackermann Property.

Keywords: Peirce's Law, modal logic, non-classical logics.

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### 1. Introduction

Peirce's Law is indeed a peculiar implicational thesis which is also of capital interest in general propositional logic. In (Salto, Méndez 1999) we have discussed the effect of Peirce's Law on some modal, intuitionistic and relevance logics. The aim of this note is to present two results posing two problems. Let us refer by PC ["Peirce classical"] and by PS5 ["Peirce S5"] to respectively, [the symbol  $\rightarrow$  will denote here strict implication]

$$(PC) \quad [(A \rightarrow B) \rightarrow A] \rightarrow A$$

and

$$(PS5) \quad [[(A \rightarrow B) \rightarrow C] \rightarrow (A \rightarrow B)] \rightarrow (A \rightarrow B)$$

Now, let  $S3P_1$  ( $S3P_2$ ) be the result of adding PS5 (PC) to Lewis' S3. We shall prove:

- (i)  $S3P_1$  and  $S3P_2$  are different logics, the first one included in the second.
- (ii)  $S3P_1$  is not included in (nor includes) Lewis' S4.
- (iii)  $S3P_1$  is included in (but does not include) Lewis' S5.
- (iv)  $S3P_2$  is not included in (nor includes) Lewis' S4 or Lewis' S5.
- (v) Both systems [ $S3P_1$  and  $S3P_2$ ] do have the Ackermann Property.

The possible interest of  $S3P_1$  and  $S3P_2$  is to be founded (among other reasons) as follows. Anderson and Belnap have reasoned at length that the Ackermann Property is necessary in any logic deserving to be called a logic of entailment (see (Anderson, Belnap 1975)). Now, though S4 and S5 do not have this property [the S4 axiom  $B \rightarrow (A \rightarrow A)$  being an immediate counterexample], it is however certainly predicabile (see §5) of both  $S3P_1$  and  $S3P_2$  [therefore, also of S3]. Then, granting as Lewis did (see (Lewis, Langford 1959)) the admissibility of paradoxes of strict implication [i.e., leaving aside questions of relevance],  $S3P_1$  and  $S3P_2$  are two strong logics of strict implication. Hence, it is natural to pose the problems:

- (vi) Is  $S3P_1$  ( $S3P_2$ ) an admissible implicational logic in Lewis' sense?
- (vii) Which are the semantics for  $S3P_1$  and  $S3P_2$ ?

## 2. Lewis' S3

Lewis' S3 can be axiomatized with (see (Lewis, Langford 1959), (Méndez 1988), (Salto, Méndez 1999)):

Axioms:

- A1.  $A \rightarrow A$
- A2.  $(A \rightarrow B) \rightarrow [(B \rightarrow C) \rightarrow (A \rightarrow C)]$
- A3.  $(B \rightarrow C) \rightarrow [(A \rightarrow B) \rightarrow (A \rightarrow C)]$
- A4.  $[A \rightarrow (A \rightarrow B)] \rightarrow (A \rightarrow B)$
- A5.  $(A \rightarrow B) \rightarrow [(C \rightarrow D) \rightarrow (A \rightarrow B)]$
- A6.  $A \rightarrow (A \vee B)$
- A7.  $B \rightarrow (A \vee B)$
- A8.  $[(A \rightarrow C) \wedge (B \rightarrow C)] \rightarrow [(A \vee B) \rightarrow C]$
- A9.  $(A \wedge B) \rightarrow A$
- A10.  $(A \wedge B) \rightarrow B$

- A11.  $[(A \rightarrow B) \wedge (A \rightarrow C)] \rightarrow [A \rightarrow (B \wedge C)]$
- A12.  $(A \rightarrow \neg A) \rightarrow \neg A$
- A13.  $(\neg A \rightarrow B) \rightarrow (\neg B \rightarrow A)$
- A14.  $A \rightarrow \neg \neg A$

Rules of derivation:

- Modus ponens: if  $\vdash A$  and  $\vdash A \rightarrow B$ , then  $\vdash B$ .
- Adjunction: if  $\vdash A$  and  $\vdash B$ , then  $\vdash A \wedge B$ .

### 3. $S3P_1$ and $S3P_2$

As noted above,  $S3P_1$  is the result of adding

$$(PS5) \quad [[(A \rightarrow B) \rightarrow C] \rightarrow (A \rightarrow B)] \rightarrow (A \rightarrow B)$$

to  $S3$ , and  $S3P_2$  the result of adding

$$(PC) \quad [(A \rightarrow B) \rightarrow A] \rightarrow A$$

to  $S3$ .

### 4. Peirce's Law and Lewis' S3

Consider the following set of matrices (I):

$\rightarrow$	0 1 2 3	$\wedge$	0 1 2 3	$\vee$	0 1 2 3	$\neg$
0	2 2 2 2	0	0 0 0 0	0	0 1 2 3	0 3
1	0 2 0 2	1	0 1 0 1	1	1 1 3 3	1 2
*2	0 0 2 2	*2	0 0 2 2	*2	2 3 2 3	*2 1
*3	0 0 0 2	*3	0 1 2 3	*3	3 3 3 3	*3 0

where designated values are starred. This set verifies  $S3P_1$  but falsifies the  $S4$  axiom  $B \rightarrow (A \rightarrow A)$  [ $v(A)=2, v(B)=3$ ] and also PC [ $v(A)=1, v(B)=0$ ].

On the other hand, the set of matrices (II), where designated values are starred:

$\rightarrow$	0 1 2 3	$\wedge$	0 1 2 3	$\vee$	0 1 2 3	$\neg$
0	2 2 2 2	0	0 0 0 0	0	0 1 2 3	0 3
*1	0 2 0 2	*1	0 2 0 1	*1	1 1 1 3	*1 2
*2	0 2 2 2	*2	0 2 2 2	*2	2 1 2 3	*2 1
*3	0 0 0 2	*3	0 1 2 3	*3	3 3 3 3	*3 0

verifies  $S3P_2$  but falsifies the  $S4$  axiom  $B \rightarrow (A \rightarrow A)$  [ $v(A)=2, v(B)=1$ ]. From these facts we immediately have:

- (a)  $S3P_1$  and  $S3P_2$  are different logics, the first one included in the second.
- (b)  $S3P_1$  does not include  $S4$ .
- (c)  $S3P_2$  does not include  $S4$ .

Now, it is known that  $S4$  is  $S3$  plus the  $S4$  axiom  $B \rightarrow (A \rightarrow A)$  and  $S5$  is  $S4$  plus  $PS5$ . It is also known that  $S3$ ,  $S4$  and  $S5$  are different logics,  $S3$  being included in  $S4$ , the latter in its turn included in  $S5$ , from which  $PC$  is not derivable. Then, we immediately have:

- (d)  $S3P_1$  does not include  $S5$ .
- (e)  $S3P_2$  does not include  $S5$ .
- (f)  $S3P_1$  is not included in  $S4$ .
- (g)  $S3P_1$  is included in  $S5$ .
- (h)  $S3P_2$  is not included in  $S4$ .
- (i)  $S3P_2$  is not included in  $S5$ .

Thus we have, as promised, propositions (i)-(v) stated in the introduction.

### 5. Ackermann Property

A logic  $L$  has the Ackermann Property (AP) just in case  $A \rightarrow (B \rightarrow C)$  is not a theorem of  $L$  when  $A$  is a propositional variable. According to Anderson and Belnap, a logic  $L$  is not a logic of implication if AP is not predicable of  $L$  (see (Anderson, Belnap 1975)). In this sense, the sets of matrices I and II verify, respectively,  $S3P_1$  and  $S3P_2$ . Now, let  $A \rightarrow (B \rightarrow C)$  be any wff where  $A$  is a wff in which neither  $\rightarrow$  nor  $\neg$  appear. Assign all variables in  $A$  the value 1. Then,  $v(A \rightarrow (B \rightarrow C))=0$  no matter the value of  $B$  and  $C$ . Thus,  $S3P_1$  and  $S3P_2$  both have the AP.

### Notes

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